

Categorical Informatics At Scale

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Conexus AI

SemWebPro 2020

$\Sigma \dashv \Delta \dashv \Pi$

Outline

The

- Who
- What
- When
- Where
- Why
- How

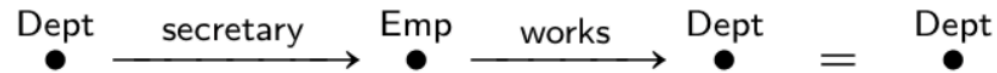
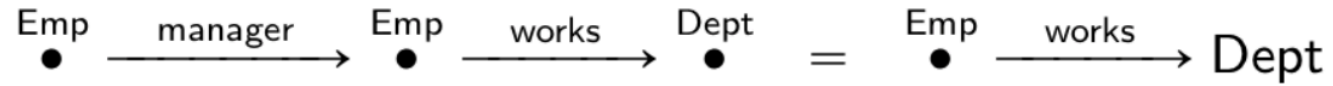
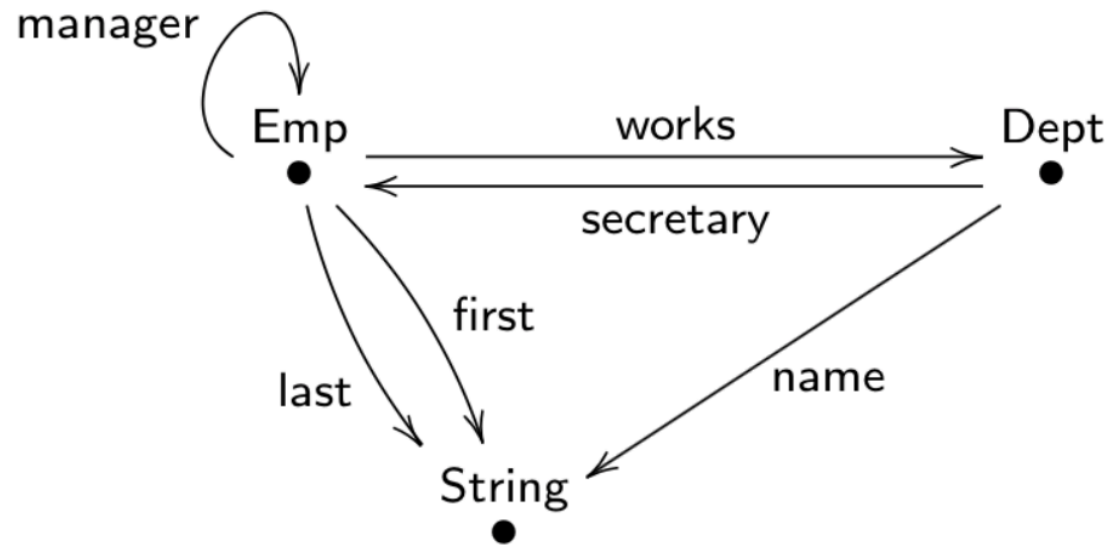
of Conexus and what it means to you, a semantic web user.

Conexus Overview & Relevance

- Who: David Spivak (math), myself (cs), Eric Daimler (business), et al.
- What: solving problems in ETL, data integration, IT interoperability, etc. (data)
- How: Kan extensions, limits and co-limits, etc. (functorial data migration)
- Why: because other technologies cannot solve these problems. (fun and profit)
- Where: San Francisco, Boston, Munich (Daniel Filonik) (also friends at NIST/DC).
- When: 2015-2018 Grant funded; 2018-present Seed stage VC funded.
- Relevance to you:
 - Use and/or contribute to the open source CQL project to do semantic web: CQL is a better OWL than OWL.
 - Use the free and/or commercial CQL IDE and/or our team to manipulate data “using real math”
 - Really hard due to non-computability issues often ignored for pragmatic reasons by mathematicians.
 - Sometimes we hire!
 - We collaborate with entrepreneurs, academics, programmers, domains experts, and more.
 - Data wrangling: it’s dirty job, but someone’s got to do it

CQL Overview

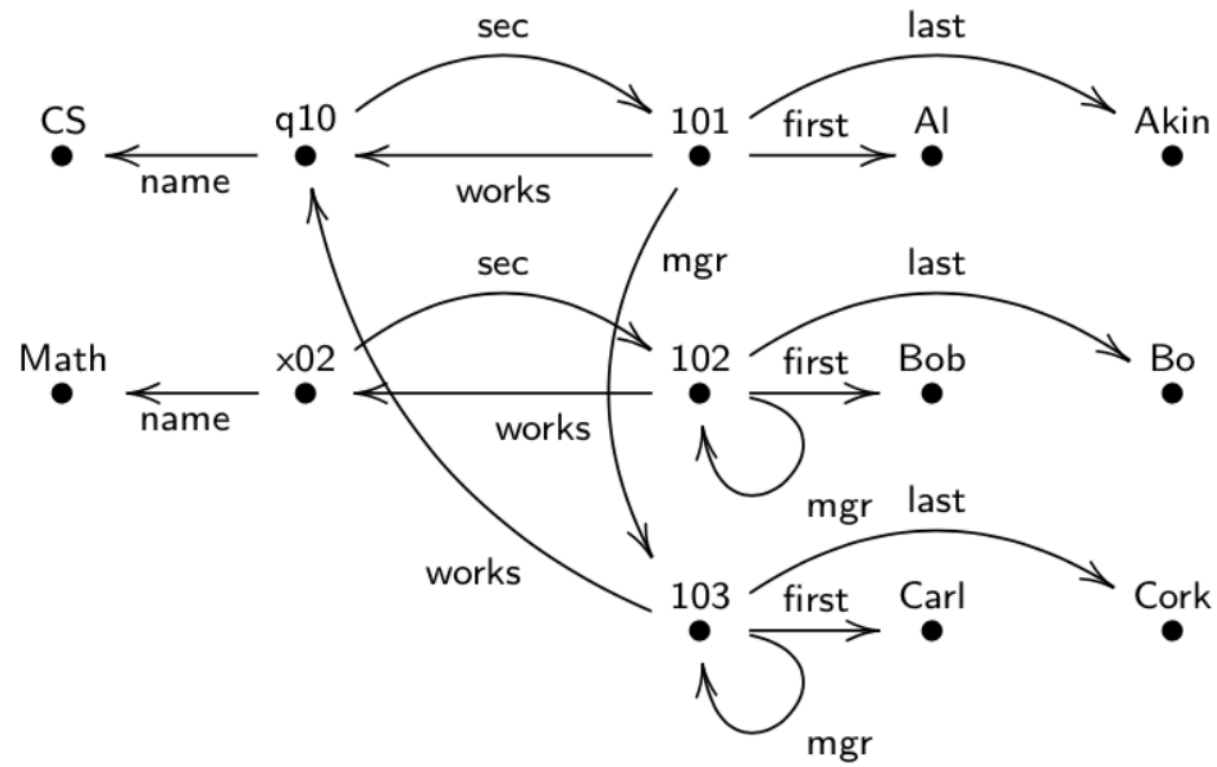
- Category theory was designed to migrate theorems from one area of mathematics to another, so it is a very natural language with which to describe how to migrate data from one schema to another.
- Community site: <http://categoricaldata.net>
- Projects:
 - NIST - several projects.
 - DARPA BRASS project.
 - Empower Retirement.
 - Stanford Chemistry Department.
 - Uber/Apache Tinkerpop (joint paper: 'algebraic property graphs')



Emp				
ID	mgr	works	first	last
101	103	q10	Al	Akin
102	102	x02	Bob	Bo
103	103	q10	Carl	Cork

Dept		
ID	sec	name
q10	101	CS
x02	102	Math

String	
ID	
Al	
Bob	
...	



Emp				
ID	mgr	works	first	last
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102	102	x02	Bob	Bo
103	103	q10	Carl	Cork

Dept		
ID	sec	name
q10	101	CS
x02	102	Math

String	
ID	
Al	
Bob	
...	

Run

New

Open

Save

Deploy

Options

<

>

Example:

Untitled 1.cql * *Employees *

```

2 ///////////////////////////////////////////////////
3
4 typeside Ty = literal {
20
21 schema S = literal : Ty {
22   entities
23     Employee
24     Department
25   foreign_keys
26     manager   : Employee -> Employee
27     worksIn   : Employee -> Department
28     secretary : Department -> Employee
29   path_equations
30     Employee.manager.worksIn = Employee.worksIn
31     Department.secretary.worksIn = Department
32   attributes
33     first last : Employee -> string
34     age       : Employee -> nat
35     cumulative_age : Employee -> nat
36     name      : Department -> string
37   observation_equations
38     forall e. cumulative_age(e) = plus(age(e), age(manager(e)))
39   options
40     prover = completion
41 }
42
43 instance I = literal : S {
65
66

```

Prove Sort

```

typeside ty
schema S
instance I : S
  equations
    first(a) = Al
    first(b) = Bob
    last(b) = Bo
    first(c) = Carl
    name(m) = Math
    name(s) = CS
    age(a) = age(c)
    manager(a) = b
    manager(b) = b
    manager(c) = c
    worksIn(a) = m
    worksIn(b) = m
    worksIn(c) = s
    secretary(s) = c
    secretary(m) = b
    secretary(worksIn(a)) = manager(a)
    worksIn(a) = worksIn(manager(a))
    age(a) = succ(succ(zero))
    age(manager(a)) = succ(zero)
  generators
    a : Employee
    b : Employee
    c : Employee
    m : Department
    s : Department

```

Employees - 12:53:27 AM

Computation wall-clock time: 0.1s
 GUI building time: 0.1s

JVM Used Change: 0 MB. Used Max: 19 MB.

Summary
 typeside Ty
 schema S
 instance I : S

Tables

TyAlg

Hom-sets

DP

Text

Expression

Department (2)

Row ▲	name	secretary
a.manager.worksIn	Math	a.manager
c.worksIn	CS	c

Employee (3)

Row ▲	age	cummulative_age	first	last	manager	worksIn
a	succ(succ(zero))	(succ(succ(zero)) plus succ(zero))	Al	a.last	a.manager	a.manager.worksIn
a.manager	succ(zero)	(succ(zero) plus succ(zero))	Bob	Bo	a.manager	a.manager.worksIn
c	succ(succ(zero))	(succ(succ(zero)) plus succ(succ(zero)))	Carl	c.last	c	c.worksIn

string (2)

Row ▲
a.last
c.last

5 IDs, 2 nulls, 0.027 seconds.

Provenance: Row limit:

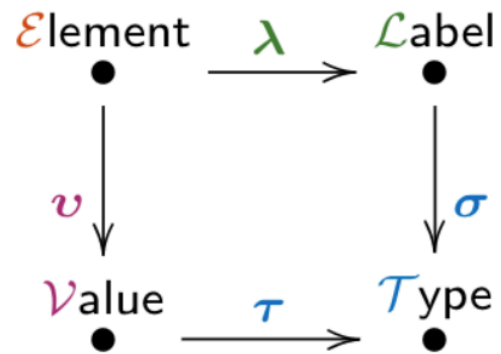
Friend of a friend

```
constraints works_at_determined1 = literal : FOAF {  
forall s:salary ->  
exists w:works_at  
where s.salary1 = w.works_at1  
      s.salary2 = w.works_at2  
}
```

```
constraints works_at_determined2 = literal : FOAF {  
forall w:works_at ->  
exists s:salary  
where s.salary1 = w.works_at1  
      s.salary2 = w.works_at2  
}
```

```
constraints frenemy_determined = literal : FOAF {  
forall fr : frenemy_of ->  
exists f : friend_of  
      e : enemy_of  
where fr.frenemy_of1 = f.friend_of1  
      fr.frenemy_of2 = f.friend_of2  
      fr.frenemy_of1 = e.enemy_of1  
      fr.frenemy_of2 = e.enemy_of2  
}
```

```
typeside Ty = literal {  
  types  
    Number  
    String  
}  
  
schema FOAF = literal : Ty {  
  entities  
    //entities  
    Person  
    Organization  
    //spans  
    knows  
    friend_of  
    works_at  
    salary  
    enemy_of  
    frenemy_of  
  foreign_keys //total functions  
    knows1 : knows -> Person  
    knows2 : knows -> Person  
    friend_of1 : friend_of -> Person  
    friend_of2 : friend_of -> Person  
    works_at1 : works_at -> Person  
    works_at2 : works_at -> Organization  
    salary1 : salary -> Person  
    salary2 : salary -> Organization  
    enemy_of1 : enemy_of -> Person  
    enemy_of2 : enemy_of -> Person  
    frenemy_of1 : frenemy_of -> Person  
    frenemy_of2 : frenemy_of -> Person  
  attributes //total functions  
    family_name : Person -> String  
    age : Person -> Number  
    given_name : Person -> String  
    salary3 : salary -> Number  
}  
  
constraints knows_symmetric = literal : FOAF {  
forall k1:knows ->  
exists k2:knows  
where k1.knows1 = k2.knows2  
      k1.knows2 = k2.knows1  
}
```



$$\begin{array}{c} \mathcal{E} \\ \bullet \end{array} \xrightarrow{v} \begin{array}{c} \mathcal{V} \\ \bullet \end{array} \xrightarrow{\tau} \begin{array}{c} \mathcal{T} \\ \bullet \end{array} = \begin{array}{c} \mathcal{E} \\ \bullet \end{array} \xrightarrow{\lambda} \begin{array}{c} \mathcal{L} \\ \bullet \end{array} \xrightarrow{\sigma} \begin{array}{c} \mathcal{T} \\ \bullet \end{array}$$

$\mathcal{E}lement$		
ID	λ	v
t_1	Trip	(u_1, u_2)
t_2	Trip	(u_1, u_3)
u_1	User	Alice
u_2	User	Bob
u_3	User	Chaz

$\mathcal{V}alue$	
ID	τ
Alice	String
Bob	String
Chaz	String
(u_1, u_2)	User \times User
(u_1, u_3)	User \times User

$\mathcal{L}abel$	
ID	σ
User	String
Trip	User \times User

$\mathcal{T}ype$
ID
String
User \times User

CQL: Further Capabilities

- CQL has over 100 keywords
- Categories, functors, and natural transformations (schema, data)
- Left and right Kan extensions (migration of data)
- Limits and colimits (integration of schema, data)
- Grothendieck construction (data \leftrightarrow schema)

- Patented implementation techniques: automated theorem proving, custom “chase” algorithms, reduction to SQL.

Thanks

- Me: ryan@conexus.com
- <http://Categoricaldata.net>
- <http://conexus.com>

- Collaborators welcome!